

Terminodes: Principles and Challenges

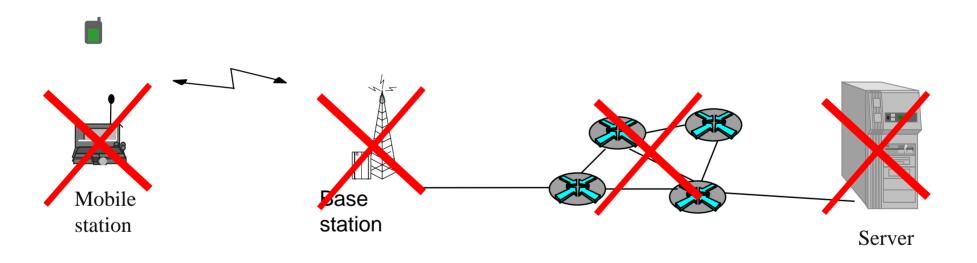
Mobile Information and Communication Systems

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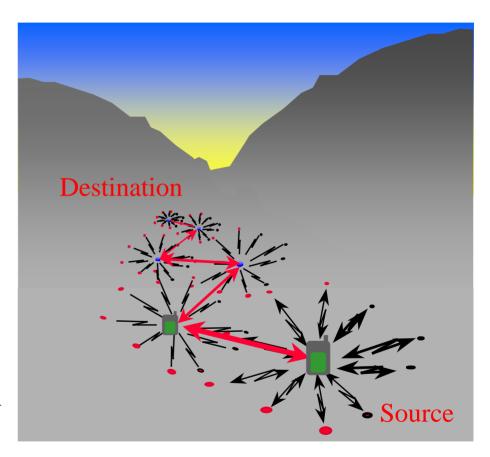
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Self-Organized, Mobile, Infrastructureless WANs



Terminal + Node = Terminode

- All **network functions** (packet forwarding, flow control, error control,...) and **terminal functions** (coding/decoding, A/D and D/A, storage, ciphering,...) are embedded in the terminode
- All terminodes are potentially **mobile**
- A communication must be **relayed** by intermediate terminodes
- The **route** followed can be different for each packet
- A terminode is able to discover its own **environment** and to react accordingly
- The network is **self-organized**: no human intervention to define the addressing plan etc.



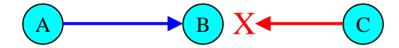
What are Terminodes?

- A societal/political vision
 - **Purpose**: make IT an instrument for democracy and economic development; empower citizens with communication facilities even in hostile environments
 - **Example**: How would infrastructureless mobile communications foster economic development in remote or poorly equipped areas?
- An intellectual fantasy
 - **Purpose**: stimulate creativity in order to identify new research challenges
 - **Example**: What would be a formal model for fair exchange?
- A technical challenge
 - **Purpose**: make innovative contributions in the area of self-organized mobile ad-hoc networks
 - **Example**: What is the best way to track the terminals in such a network?

Reminder: Packet Radio

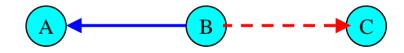
- Research started in the 70's, essentially with military applications in mind
- Unlike with cellular networks, the allocation of the transmission resource is decentralized (no base stations, no cells)
- Two typical problems:

The hidden terminal



- A is sending to B
- C is out of the range of A's transmitter
- C wants to send to B (or someone near to B); a collision occurs in B
- A is *hidden* from C

The exposed terminal



- B is sending to A
- C is in the range of B's transmitter
- C wants to send, but will wait; if A is out of the range of C, then C waits needlessly
- C is *exposed* to B

Upper Bound for the Throughput of Packet Radio Networks

If we have:

- *n* identical randomly located nodes
- each capable of transmitting W bits/s Then the throughput $\lambda(n)$ obtainable by each node for a *randomly chosen* destination is

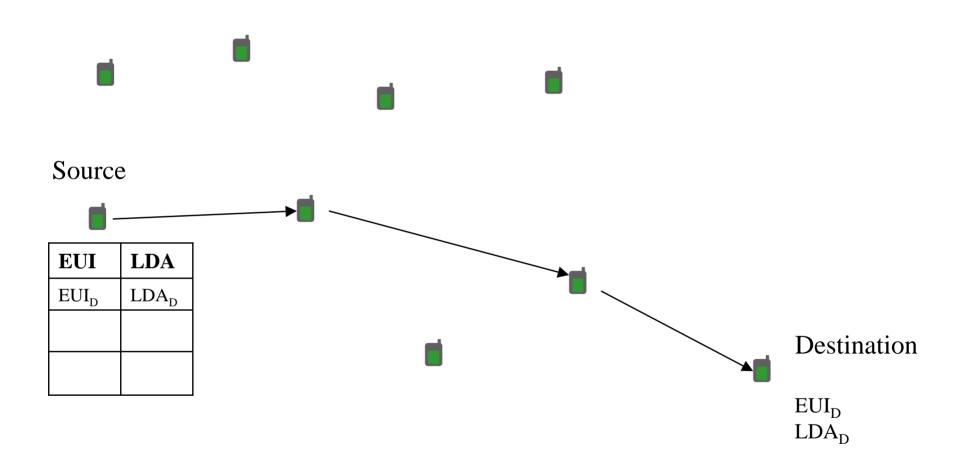
$$\lambda(n) = \Theta\left(\frac{W}{\sqrt{n\log n}}\right)$$

Ref: P. Gupta, P. Kumar, *The Capacity of Wireless Networks* IEEE Transactions on Information Theory, March 2000

Technical Issues covered by this Talk

- Basic mechanism for terminodes: packet forwarding, location awareness
- Mobility management of terminodes
- Incentive to collaborate for terminodes

Packet Forwarding in Terminodes

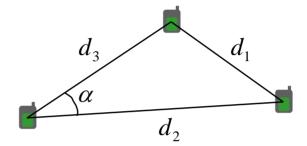


EUI: End-system Unique Identifier (64 bits)

LDA: Location-Dependent Address: (longitude, latitude, height) (approx. 48 bits for 10m accuracy)

Location Awareness





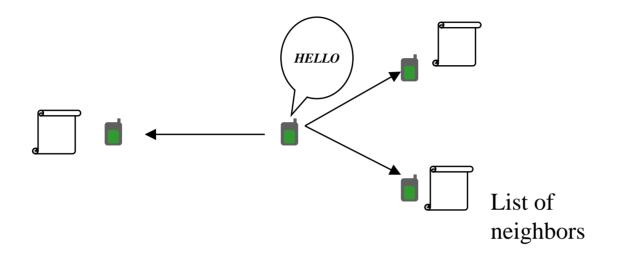
$$\cos \alpha = \frac{d_1^2 - d_2^2 - d_3^2}{2d_2d_3}$$

If GPS is available: Each node is aware of its own Location-Dependent Address (LDA) via GPS (Global Positioning System).

If GPS is not available: computation of relative positions based on Time of Arrival.

Major Pb: Non-Line-of-Sight

Neighborhood Awareness

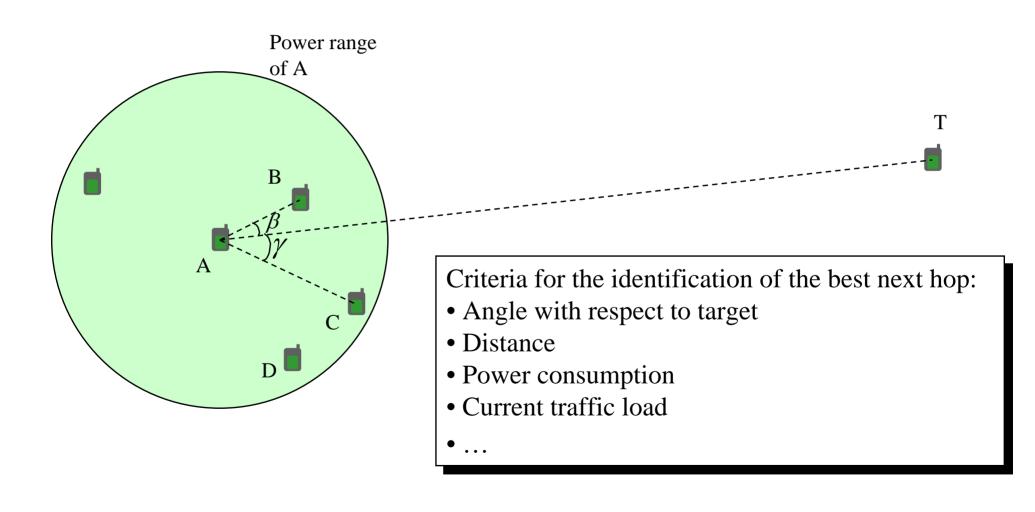


Periodic Broadcast of *Hello* Messages with current position of the sender.

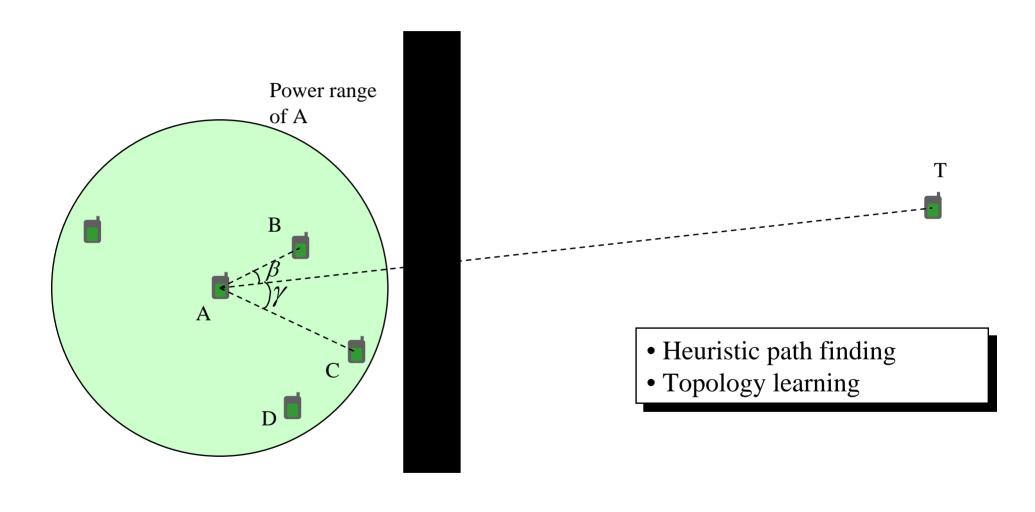
Applications:

- Code negotiation (for spread spectrum)
- Power control
- Establishment of pairwise security keys

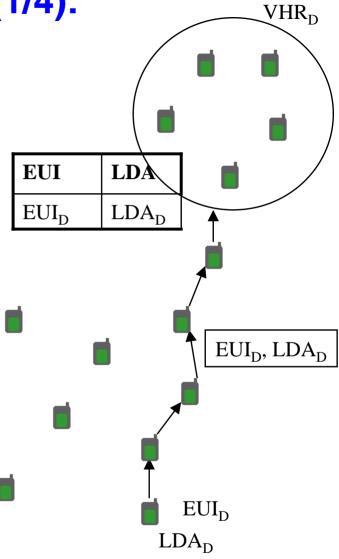
Directional Packet Forwarding



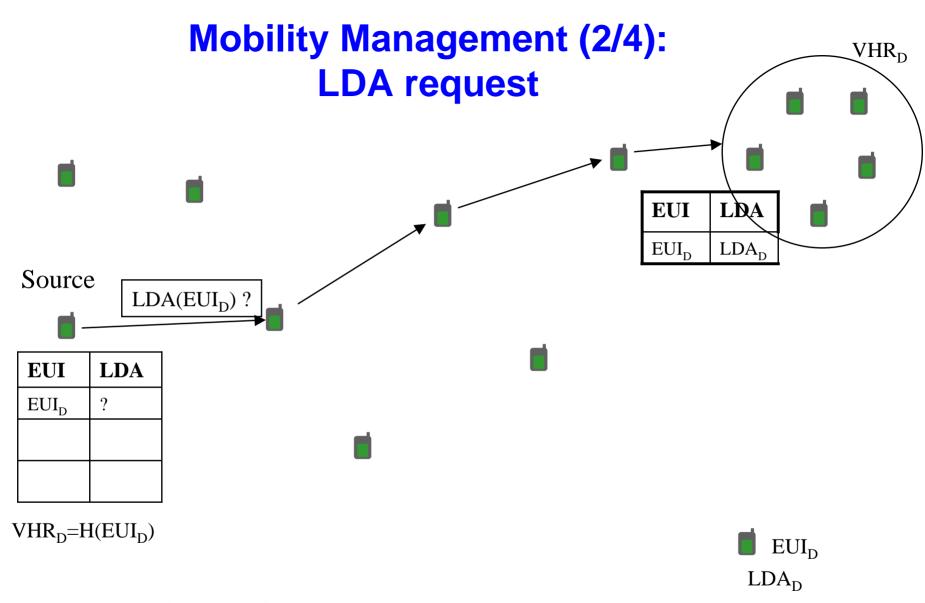
Obstacle Avoidance



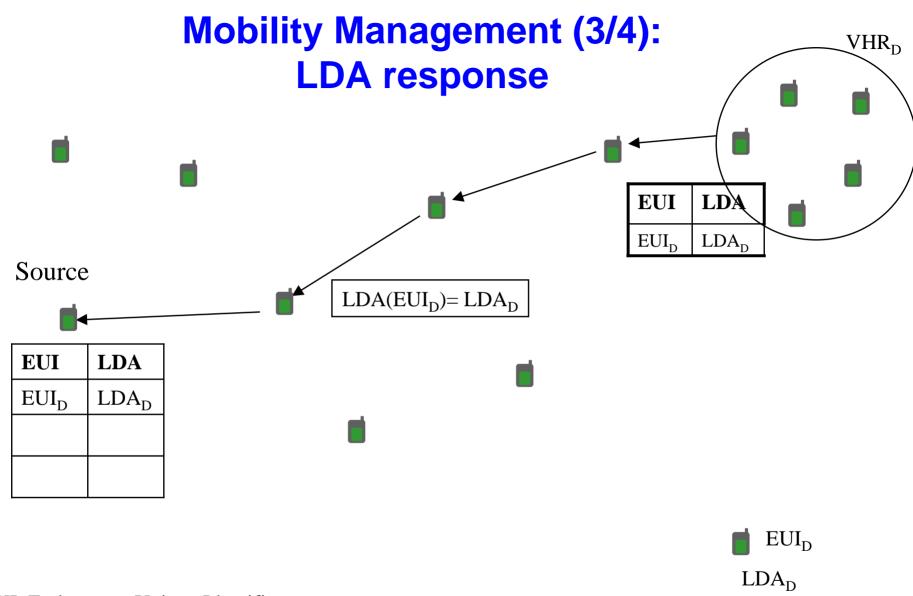
Mobility Management (1/4): LDA storage



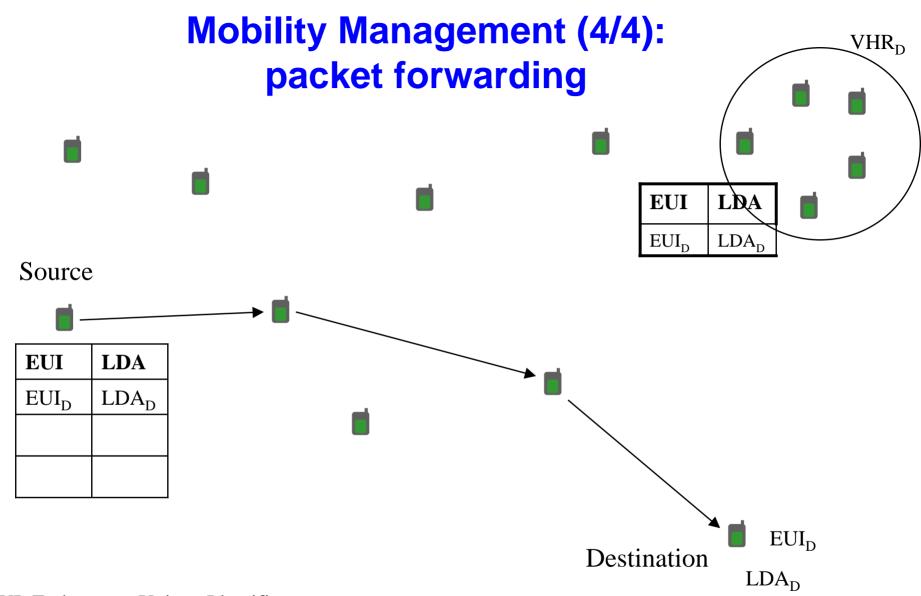
EUI: End-system Unique Identifier LDA: Location-Dependent Address



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Incentive to cooperate and to prevent congestion

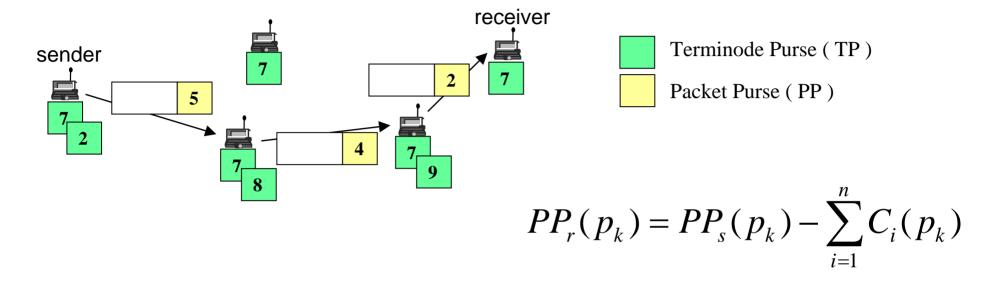
• Mechanism required to:

- Encourage end-users to let their terminode act as a relay (keep them turned on and not tamper with them)
- Discourage end-users from **overloading** the network; in particular, limit the number of long distance communications

• 2 models

- Packet Purse Model (payment by the sender)
- Packet Trade Model (payment by the receiver)

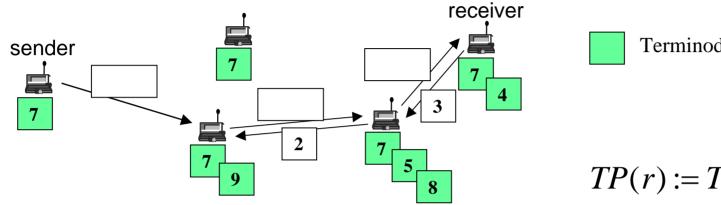
Charging the sender: Packet Purse Model (PPM)



Problems to be solved include:

- Nugget forgery should be prevented
- Nugget robbery should be prevented
- Each packet should **indeed** be forwarded
- The packet purse should be **bundled** to its packet

Charging the receiver: Packet Trade Model (PTM)



Terminode Purse (TP)

$$TP(r) := TP(r) - \sum_{i=1}^{n} C_i(p_k)$$

Advantages

- The sender does not need to know the **amount of nuggets** that is necessary to send a packet
- Intermediaries are interested in forwarding the packet after having bought it
- Charging for **multicast** communications is easier

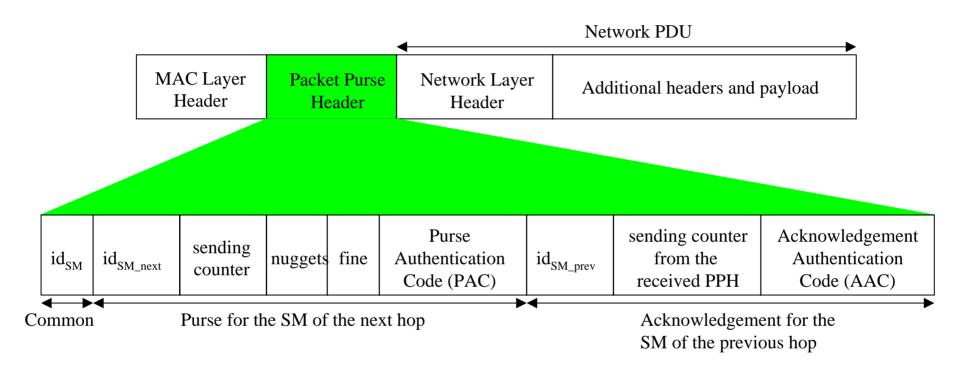
Drawback

There is no direct incentive to refrain from overloading the network

Implementation: Assumptions

- Tamper resistant security module in each device, which is used for the management of nuggets and cryptographic keys
- Public key infrastructure that can be used by the security modules to authenticate each other and establish secure communication links
- Neighborhood changes slowly (at least compared to the speed of the packets)
- Reliable, bidirectional communication between neighbors
- Pricing mechanism
- Terminodes are **greedy** and they always want to increase their number of nuggets

Implementation of PPM: Packet Purse Header (PPH)

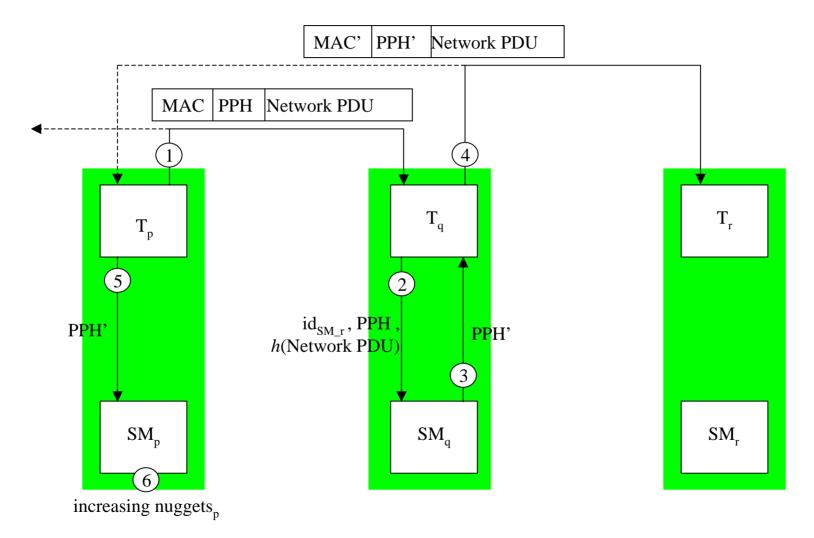


 $PAC = g(k_{SM,SM_next}; id_{SM}, id_{SM_next}, sending counter, nuggets, fine, h(Network PDU))$

 $AAC = g(k_{SM,SM_prev}; received PPH)$

g is a publicly known keyed cryptographic hash function h is a publicly known cryptographic hash function

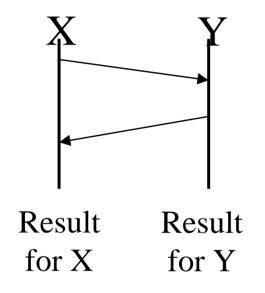
Implementation of PPM: Packet forwarding protocol



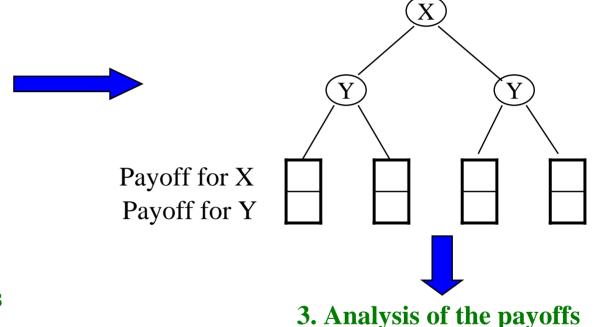
PPH: Packet Purse Header

Modeling exchange protocols using game theory

1. Protocol Design



2. Modelling as a game tree

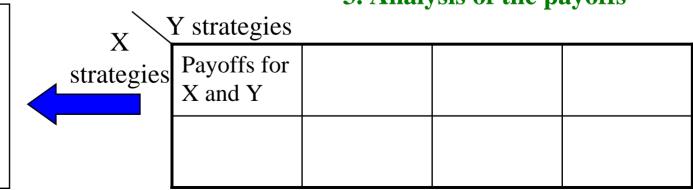


4. Identification of the properties

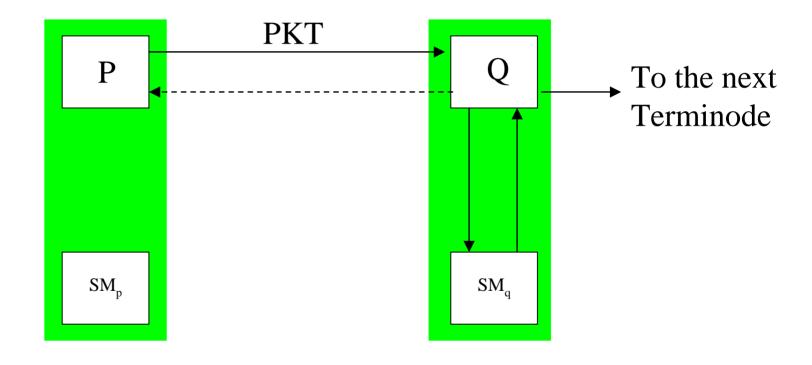
- Fairness
- Incentive for desirable behaviours

Example of practical usage:

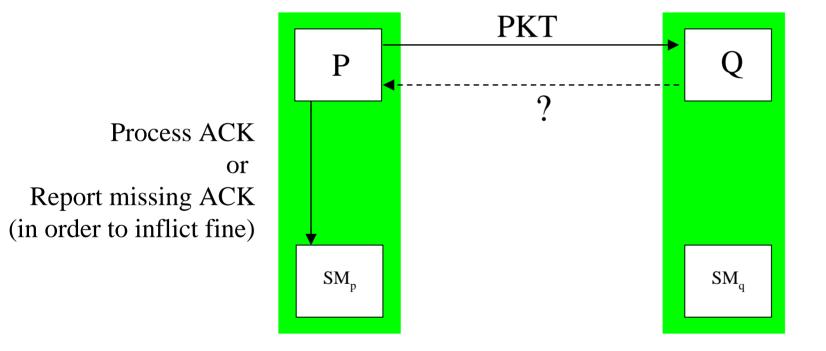
• Fine tuning of the parameters



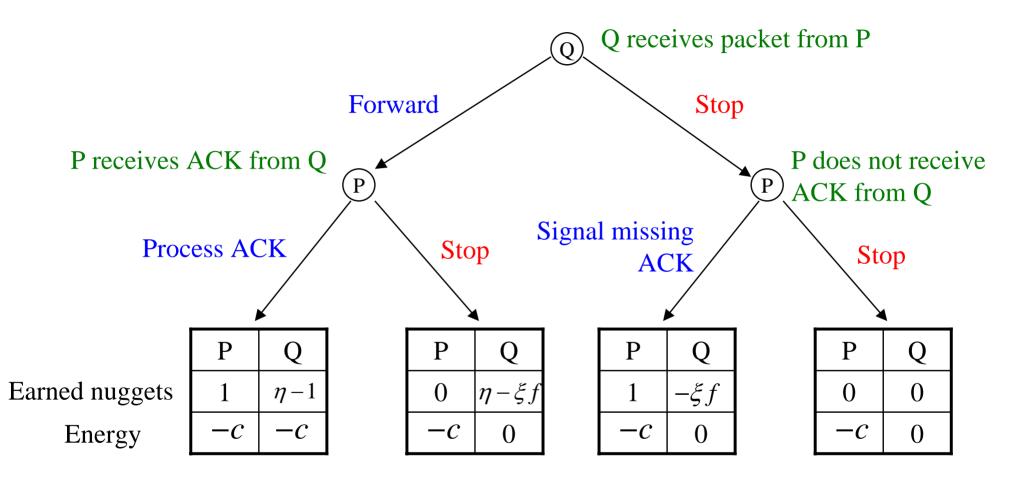
Modeling packet forwarding as a game: Q's behavior



Modeling packet forwarding as a game: P's behavior



Payoff for each player in packet forwarding



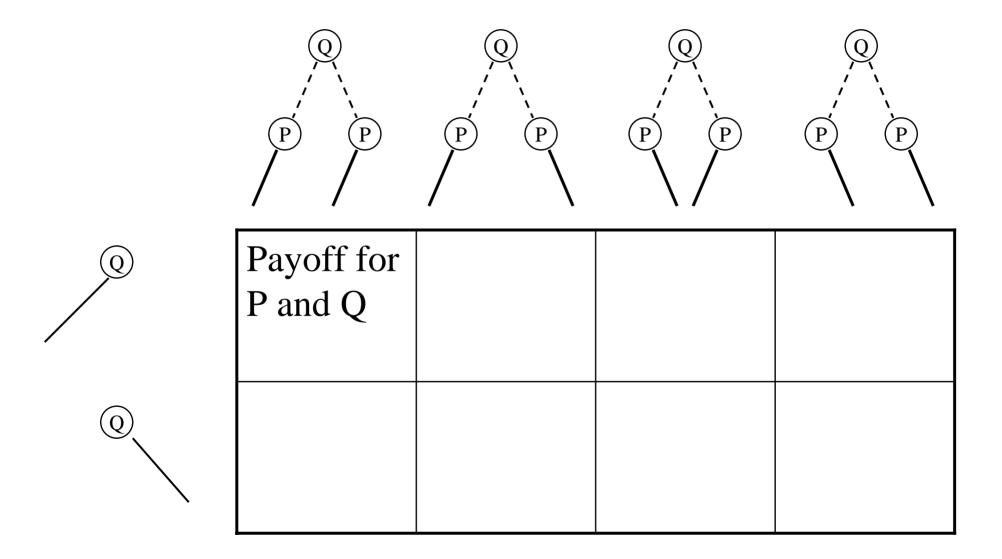
C: cost

f: fine

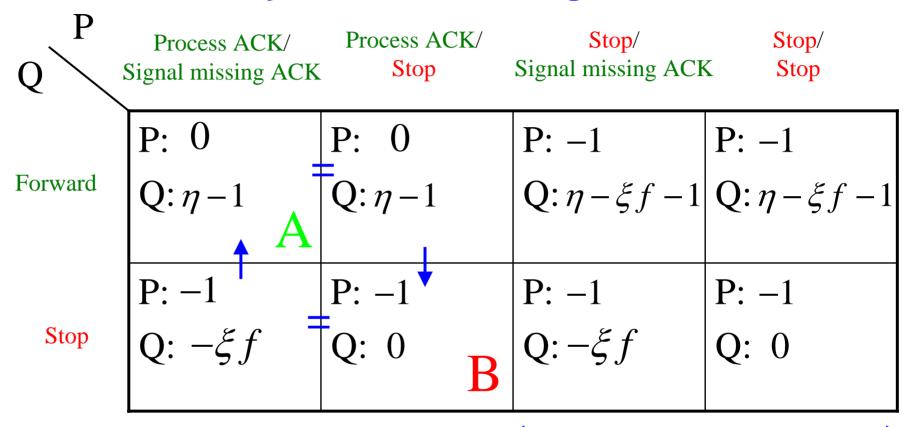
 η : probability that Q will eventually receive payment from R

 ξ : probability that the fine will eventually be inflicted to Q

Identification of the strategies



Analysis of the strategies



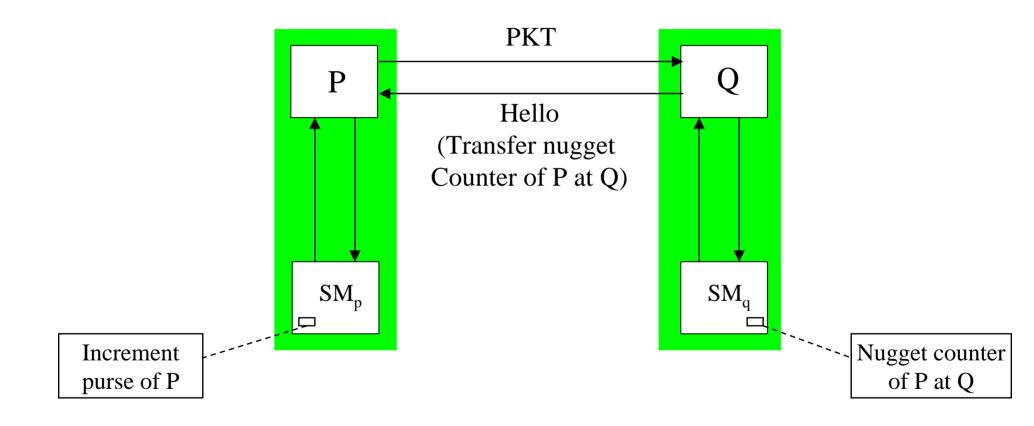
$$\eta - 1 > -\xi f \Rightarrow f > \frac{1 - \eta}{\xi}$$
 (1)

Weakly dominated by the two left columns

A: desired Nash equilibrium

B: undesired Nash equilibrium

Alternative solution (non real-time increase of the Terminode Purse)



Conclusion on the incentive to cooperation

- In self-organized mobile ad hoc networks, the **cooperation** mechanisms have to be carefully scrutinized
- Some form of **funny money** (nuggets) can be introduced to incentivate cooperation
- Game theory can be used to study the properties of exchange protocols

Future work:

- Reduce the overhead (statistical mechanisms?)
- Get rid of the assumption of communication bidirectionality

Main challenge and benefit: Working accross layers

Mathematical foundation	Information theory	Security	Economics	System architectu	Communicating embedded systems
			Real-time services		
			Information systems		
			Network layer		
ation	y		Physical layer and MAC	re	1S

Related Projects

- Carnegie-Mellon: Monarch
 - Testbed
 - Routing protocols (proposal: Dynamic Source Routing Protocol)
- MIT/LCS: Wireless Network of Devices (WIND)
 - Energy-efficient routing
 - Intentional Naming System
- Cornell Univ./EE Dept
 - Zone Routing Protocol (ZRP)
 - Mobility management with Uniform Quorum Systems (UQS) and Randomized Database Group (RDG)
- Georgia Tech/School of ECE: Associativity-Based Routing (ABR)
- SUN: Ad Hoc on Demand Distance Vector Routing (AODV)
- UCLA/CS Dept: Clustering
- Univ. of Texas at Dallas/CS Dept: Clustering
- IETF: MANET Working Group
- ...

References on the Project

- Terminodes: Toward Self-Organized Mobile Wide Area Networks JP Hubaux, Technical Report SSC/1999/022, June 3, 1999
- The Terminode Project: Towards Mobile Ad-Hoc WANs
 JP Hubaux, JY Le Boudec, S. Giordano, M. Hamdi,
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- Toward Mobile Ad-Hoc WANs: Terminodes
 JP Hubaux, JY Le Boudec, M. Vojnovic, S. Giordano, M. Hamdi, L. Blazevic, L. Buttyan
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 Technical Report SSC/2000/025, April 2000; accepted for publication at MobiHoc'00
- Toward a Formal Model of Fair Exchange a Game Theoretic Approach
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 Technical report No. SSC/1999/039, December 1999

All documents available at http://dscwww.epfl.ch. See also http://www.terminodes.org